

P197

$$\begin{aligned}
 18. \sum_{n=1}^{\infty} n p_n - \sum_{n=1}^{\infty} r_n &= \sum_{n=1}^{\infty} n(r_n - r_{n+1}) - \sum_{n=1}^{\infty} r_n \\
 &= \sum_{n=1}^{\infty} (n-1)r_n - \sum_{n=1}^{\infty} n r_{n+1} \\
 &= (1-1)r_1 + \sum_{n=2}^{\infty} (n-1)r_n - \sum_{n=1}^{\infty} n r_{n+1} \\
 &= \sum_{n=1}^{\infty} n r_{n+1} - \sum_{n=1}^{\infty} n r_{n+1} \\
 &= 0
 \end{aligned}$$

故 $\sum_{n=1}^{\infty} n p_n = \sum_{n=1}^{\infty} r_n$

$$E(X) = \sum_{n=1}^{\infty} n p_n = \sum_{n=1}^{\infty} r_n = \sum_{n=1}^{\infty} P(X \geq n)$$

20. 若 $Y = \min\{X_1, X_2, \dots, X_m\}$

$$E(Y) = \sum_{n=1}^{\infty} P(Y \geq n)$$

$$P(Y \geq n) = P(X_1 \geq n, X_2 \geq n, \dots, X_m \geq n)$$

由于 X_1, X_2, \dots, X_m 是独立的。

$$\text{故 } P(Y \geq n) = P(X_1 \geq n) P(X_2 \geq n) \cdots P(X_m \geq n) = r_n \cdot r_n \cdots r_n = r_n^m$$

$$\text{故 } E(Y) = \sum_{n=1}^{\infty} r_n^m$$

$$21. F_X(u) = \int_{-\infty}^u f(t) dt$$

$$F_X(u) + r(u) = \int_{-\infty}^{+\infty} f(t) dt = 1, \text{ 故 } r(u) = 1 - F_X(u)$$

$$\therefore F_X(u) = P(X \leq u)$$

$$\therefore r(u) = 1 - P(X \leq u) = P(X \geq u)$$

$$\begin{aligned}
 \text{故 } \int_0^{\infty} P(X \geq u) du &= \int_0^{+\infty} r(u) du = \int_0^{+\infty} \underbrace{1 - F_X(u)}_{u=0} du \Big|_{u=0}^{u \rightarrow +\infty} \\
 &= \lim_{A \rightarrow +\infty} [r(u)] \Big|_{u=0}^{u=A} = \int_0^A r(u) du
 \end{aligned}$$

$$= 0 + \int_0^{+\infty} u f(u) du$$

$$= \int_0^{+\infty} u f(u) du$$

$$\text{由于 } X \text{ 非负, 故 } E(X) = \int_{-\infty}^{+\infty} u f(u) du = \int_0^{+\infty} u f(u) du$$



$$22. f(t) = \lambda e^{-\lambda t} \quad (t \geq 0)$$

$$E(X) = \int_0^{+\infty} t f(t) dt = \int_0^{+\infty} \lambda t e^{-\lambda t} dt = \left[\lambda t - \left(t + \frac{1}{\lambda} \right) e^{-\lambda t} \right] \Big|_0^{+\infty} = \frac{1}{\lambda}$$

$$r(u) = \int_u^{+\infty} f(t) dt = \int_u^{+\infty} \lambda e^{-\lambda t} dt = e^{-\lambda u} \quad (u \geq 0)$$

$$r(u) = \begin{cases} e^{-\lambda u}, & u \geq 0 \\ 1, & u < 0 \end{cases}$$

$$\int_0^{+\infty} r(u) du = \int_0^{+\infty} e^{-\lambda u} du = -\frac{1}{\lambda} e^{-\lambda u} \Big|_0^{+\infty} = \frac{1}{\lambda}$$

$$\frac{1}{\lambda} E(X) = \int_0^{+\infty} r(u) du = \frac{1}{\lambda}$$

23.

$$\begin{aligned} E(T) &= \int_0^{+\infty} P(T \geq t) dt = \alpha \int_0^{+\infty} e^{-\lambda t} dt + (1-\alpha) \int_0^{+\infty} e^{-\mu t} dt \\ &= \frac{\alpha}{\lambda} + \frac{1-\alpha}{\mu} \end{aligned}$$

$$\begin{aligned} P(T^2 \geq t) &= P(T \geq \sqrt{t}) + P(T \leq -\sqrt{t}) \quad P(T \leq -\sqrt{t}) = 0 \\ &= \alpha e^{-\lambda \sqrt{t}} + (1-\alpha) e^{-\mu \sqrt{t}} \end{aligned}$$

$$\frac{1}{2} E(T^2) = \int_0^{+\infty} P(T^2 \geq t) dt = \int_0^{+\infty} \alpha e^{-\lambda \sqrt{t}} + (1-\alpha) e^{-\mu \sqrt{t}} dt$$

$$\stackrel{s=\sqrt{t}}{\Rightarrow} s = \sqrt{t}, \quad ds = \frac{1}{2\sqrt{t}} dt = \frac{1}{2s} dt, \quad dt = 2sds$$

$$\begin{aligned} LHS &\Rightarrow 2\alpha \int_0^{+\infty} s e^{-\lambda s} ds + 2(1-\alpha) \int_0^{+\infty} s e^{-\mu s} ds \\ &= 2\alpha \left(-\frac{\lambda s+1}{\lambda^2} e^{-\lambda s} \right) \Big|_0^{+\infty} + 2(1-\alpha) \left(-\frac{\mu s+1}{\mu^2} e^{-\mu s} \right) \Big|_0^{+\infty} \\ &= \frac{2\alpha}{\lambda^2} + \frac{2(1-\alpha)}{\mu^2} \end{aligned}$$

$$\sigma^2(T) = E(T^2) - [E(T)]^2 = \frac{2\alpha}{\lambda^2} + \frac{2(1-\alpha)}{\mu^2} - \left(\frac{\alpha}{\lambda} + \frac{1-\alpha}{\mu} \right)^2$$

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24. 设可以再用 Y 小时, $Y \in [0, +\infty)$, 已经用了 X 小时

$$P(Y \geq u) = \frac{P(Y \geq u, X \geq n)}{P(X \geq n)} \quad f(t) = e^{-\lambda t} \text{ 是 } X \text{ 的密度函数}$$

$$P(Y \geq u | X \geq n) = \frac{P(Y \geq u, X \geq n)}{P(X \geq n)} = \frac{P(X \geq u+n)}{P(X \geq n)} = \frac{\int_{u+n}^{+\infty} \lambda e^{-\lambda t} dt}{\int_n^{+\infty} \lambda e^{-\lambda t} dt} = \frac{e^{-\lambda(u+n)}}{e^{-\lambda n}} = e^{-\lambda u}$$

$$E(Y | X \geq n) = \int_0^{+\infty} P(Y \geq u | X \geq n) du$$

$$= \int_0^{+\infty} e^{-\lambda u} du$$

$$= \frac{1}{\lambda}$$

无论是否已经用了多长时间, 接着能用时间的期望总是 $\frac{1}{\lambda}$

29. 记每次选到的点分别为 X_1, X_2, \dots, X_n , $X_j \in [0, 1]$, $1 \leq j \leq n$

$$Y = \max \{X_1, \dots, X_n\}, \quad Z = \min \{X_1, \dots, X_n\}$$

X_j 的生成函数 $f_{X_j}(u) = u$, $1 \leq j \leq n$, $0 \leq u \leq 1$

$$F_Y(u) = P(Y \leq u) = P(X_1 \leq u, X_2 \leq u, \dots, X_n \leq u)$$

由于 X_1, \dots, X_n 是独立的

$$\text{故 } F_Y(u) = [P(X_j \leq u)]^n = u^n \quad 0 \leq u \leq 1$$

$$f_Y(u) = F'_Y(u) = n u^{n-1}$$

$$E(Y) = \int_0^1 u f_Y(u) du = \int_0^1 n u^n du = \frac{n}{n+1} u^{n+1} \Big|_0^1 = \frac{n}{n+1}$$

$$P(Z \geq u) = P(X_1 \geq u, X_2 \geq u, \dots, X_n \geq u) = [P(X_j \geq u)]^n = (1-u)^n$$

$$E(Z) = \int_0^1 P(Z \geq u) du = \int_0^1 (1-u)^n du = \int_0^1 t^n dt = \frac{1}{n+1} t^{n+1} \Big|_0^1 = \frac{1}{n+1}$$

$$\text{区间长度 } T = Y - Z, \text{ 故 } E(T) = E(Y) - E(Z) = \frac{n}{n+1}$$



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33. 对 $\forall 1 \leq j \leq n$, X_j 的生成函数 $g_j(z) = \frac{1}{4}(z^{-1} + \frac{1}{2} + \frac{1}{4}z)$

$\{X_j\}$ 相互独立，则 S_n 生成函数 $g(z) = g_1(z) \cdots \cdot g_n(z) = (\frac{1}{4}z^{-1} + \frac{1}{2} + \frac{1}{4}z)^n$

$P(S_n=0)$ 是 $g(z)$ 中常数项

$$g(z) = \left(\frac{1}{4}z^{-1} + \frac{1}{2} + \frac{1}{4}z\right)^n = \left(\frac{1+2z+z^2}{4z}\right)^n = \frac{(z+1)^{2n}}{4^n z^n}$$

$(z+1)^{2n}$ 中 z^n 前系数 $\binom{2n}{n}$

$$\therefore P(S_n=0) = \frac{\binom{2n}{n}}{4^n}$$

设 A, B 掷硬币掷出正面次数 Y_1, Y_2

$$P(Y_j=k) = \binom{n}{k} \left(\frac{1}{2}\right)^k \left(\frac{1}{2}\right)^{n-k} = \frac{\binom{n}{k}}{2^n} \quad \text{for } j=1 \text{ or } 2$$

Y_1 生成函数 $h_1(z) = \sum_{k=0}^n \frac{\binom{n}{k}}{2^n} z^k = \frac{(z+1)^n}{2^n}$

$-Y_2$ 生成函数 $h_2(z) = \sum_{k=0}^n \frac{\binom{n}{k}}{2^n} z^{-k} = \frac{(\frac{1}{z}+1)^n}{2^n}$

令 $Z = Y_1 - Y_2$, Z 生成函数 $h(z) = h_1(z)h_2(z) = \frac{(1+z)^{2n}}{4^n z^n}$
常数项为 $\frac{\binom{2n}{n}}{4^n}$, 代表 $P(Z=0)$ 即 $P(Y_1=Y_2)$ 的概率

$$37. (1) \int_{-\infty}^{+\infty} e^{-\lambda u} f(u) du = \int_0^c \frac{1}{c} e^{-\lambda u} du = \frac{1}{c\lambda} e^{-\lambda u} \Big|_0^c = \frac{1 - e^{-\lambda c}}{\lambda c} \quad (\lambda > 0)$$

$$(2) \int_{-\infty}^{+\infty} e^{-\lambda u} f(u) du = \int_0^c \frac{2ue^{-\lambda u}}{c^2} du = \frac{2(u+1)e^{-\lambda u}}{c^2 \lambda^2} \Big|_0^c = \frac{2[1 - (c+1)e^{-\lambda c}]}{\lambda c^2} \quad (\lambda > 0)$$

$$(3) \int_{-\infty}^{+\infty} f(u) du = \int_0^{+\infty} \frac{\lambda^n u^{n-1}}{(n-1)!} e^{-\lambda u} du$$

$$= \lim_{A \rightarrow \infty} -\frac{\lambda^n u^{n-1}}{(n-1)!} e^{-\lambda u} \Big|_{u=0}^{u=A} + \int_0^{+\infty} \frac{\lambda^n u^{n-1}}{(n-1)!} e^{-\lambda u} du$$

$$= \int_0^{+\infty} \frac{\lambda^n u^{n-1}}{(n-1)!} e^{-\lambda u} du$$

$$\therefore I_n = \int_0^{+\infty} \frac{\lambda^{n-1} u^{n-1}}{(n-1)!} e^{-\lambda u} du$$

故 $\int_{-\infty}^{+\infty} f(u) du = 1$, $f(u)$ 是密度函数



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$$\Gamma(s) = \int_0^{+\infty} e^{-su} f(u) du = \int_0^{+\infty} \frac{\lambda^n u^{n-1}}{(n-1)!} e^{-(\lambda+s)u} du$$

$$= \frac{\lambda^n}{(\lambda+s)^n} \int_0^{+\infty} \frac{(\lambda+s)^n u^{n-1}}{(n-1)!} e^{-(\lambda+s)u} du, \text{ 积分是密度的积分, 结果为 } \\ = \frac{\lambda^n}{(\lambda+s)^n}$$

38. T_i 的密度 $f_{T_i}(t) = \lambda e^{-\lambda t}$

$$\int_0^{+\infty} e^{-st} f_{T_i}(t) dt = \lambda \int_0^{+\infty} e^{-(s+\lambda)t} dt \\ = \frac{\lambda}{s+\lambda} \int_0^{+\infty} e^{-(s+\lambda)t} dt, \text{ 积分是密度的积分, 结果为 } 1$$

$$= \frac{\lambda}{s+\lambda}$$

S_n 的生成函数的 T_i 的乘积

故 S_n 的拉普拉斯变换 $\Gamma(s) = \frac{\lambda^n}{(s+\lambda)^n}$, 与上题相同



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