

吴诗非 2020/03/89

$$11. P(A|C) = \frac{P(AC)}{P(C)}, \quad P(B|C) = \frac{P(BC)}{P(C)}$$

$$P(A|C) \geq P(B|C) \Rightarrow \frac{P(AC)}{P(C)} \geq \frac{P(BC)}{P(C)}$$

$$\therefore P(C) > 0$$

$$\therefore P(AC) \geq P(BC)$$

$$\text{同理}, P(AC^c) \geq P(BC^c) \quad (1),$$

$$(1) + (2), P(AC) + P(AC^c) \geq P(BC) + P(BC^c)$$

$$P(AC) + P(AC^c) = P(AC \cup AC^c) = P(A \cap \Omega) = P(A)$$

$$P(BC) + P(BC^c) = P(B)$$

$$\therefore P(A) \geq P(B)$$

$$16. (1) \text{ 只需证 } P((A_1 \cup A_2)A_3 \cap (A_4^c \cup A_5^c)) = P((A_1 \cup A_2)A_3)P(A_4^c \cup A_5^c) \text{ 成立}$$

$$\text{左边} = P((A_1 \cup A_2) \cap (A_3(A_4^c \cup A_5^c))) = P(A_1 A_3 A_4^c \cup A_1 A_3 A_5^c)$$

$$= P(A_1 A_3 (A_4^c A_5^c)) + P(A_2 A_3 (A_4^c A_5^c)) - P(A_1 A_2 A_3 (A_4^c A_5^c))$$

$$= P(A_1 A_3 A_4^c) + P(A_1 A_3 A_5^c) - P(A_1 A_2 A_3 A_4^c A_5^c) + P(A_2 A_3 A_4^c) + P(A_2 A_3 A_5^c) - P(A_2 A_3 A_4^c A_5^c)$$

$$= P_1 P_3 (1-P_4) + P_1 P_3 (1-P_5) + P_1 P_3 (1-P_4) (1-P_5) + P_2 P_3 (1-P_4) + P_2 P_3 (1-P_5) - \dots$$

$$\text{右边} = [P(A_1 A_2) + P(A_2 A_3) - P(A_1 A_2 A_3)] [1 - P(A_4^c A_5^c)]$$

$$= (P_1 + P_2 - P_1 P_2) P_3 (1 - P_4 P_5) = \text{左边}$$

故 $(A_1 \cup A_2)A_3$ 与 $A_4^c \cup A_5^c$ 独立

$$(2) P((A_1 \cup A_2) \cap (A_3 \cap A_4)) = P(A_1 A_3 A_4 \cup A_2 A_3 A_4) = P(A_1 A_3 A_4) + P(A_2 A_3 A_4) - P(A_1 A_2 A_3 A_4)$$

$$= (P_1 + P_2 - P_1 P_2) P_3 P_4$$

$$= P(A_1 \cup A_2) P(A_3 A_4)$$

$$P(A_3 \cap A_4) \cap A_5^c = P(A_3 A_4 - A_3 A_4 A_5) = P(A_3 A_4) - P(A_3 A_4 A_5) = P(A_3 A_4) [1 - P(A_5)] = P(A_3 A_4) P(A_5^c)$$

$$P(A_5^c \cap (A_1 \cup A_2)) = P(A_1 A_5^c) + P(A_2 A_5^c) - P(A_1 A_2 A_5^c)$$

$$= [P(A_1) + P(A_2) - P(A_1) P(A_2)] P(A_5^c)$$

$$= P(A_1 \cup A_2) P(A_5^c)$$

$$P((A_1 \cup A_2) \cap A_3 A_4 \cap A_5^c) = P(A_1 A_3 A_4 A_5^c) + P(A_2 A_3 A_4 A_5^c) - P(A_1 A_2 A_3 A_4 A_5^c)$$

$$= [P(A_1) + P(A_2) - P(A_1 A_2)] P(A_3 A_4) P(A_5^c)$$

$$= P(A_1 \cup A_2) P(A_3 A_4) P(A_5^c)$$

故 $A_1 \cup A_2, A_3 \cap A_4, A_5^c$ 相互独立



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17. 设 $A_i = \{\text{该老虎机是第}i\text{个老虎机}\}$, $P(A_1) = P(A_2) = P(A_3)$
 $B = \{\text{4次中获得2次回报}\}$

$$P(A_1) = \frac{P(B|A_1)}{P(B|A_1) + P(B|\bar{A}_1)} \quad P(B) = P(A_1) \cdot C_4^2 \left(\frac{1}{3}\right)^2 \left(\frac{2}{3}\right)^2 + P(A_2) \cdot C_4^1 \left(\frac{1}{2}\right)^1 \left(\frac{1}{2}\right)^3 + P(A_3) \cdot C_4^1 \left(\frac{2}{3}\right)^1 \left(\frac{1}{3}\right)^3 \\ = \frac{16}{81} + \frac{1}{8} = \frac{209}{648}$$

$$P(A_1|B) = \frac{P(A_1) \cdot C_4^2 \left(\frac{1}{3}\right)^2 \left(\frac{2}{3}\right)^2}{P(B)} \quad P(A_1|B) = \frac{P(A_1|B)}{P(B)} = \frac{P(A_1) \cdot C_4^2 \left(\frac{1}{3}\right)^2 \left(\frac{2}{3}\right)^2}{P(B)} = \frac{64}{209} \\ P(A_2|B) = \frac{P(A_2|B)}{P(B)} = \frac{P(A_2) \cdot C_4^1 \left(\frac{1}{2}\right)^1 \left(\frac{1}{2}\right)^3}{P(B)} = \frac{1}{8} = \frac{81}{209} \\ P(A_3|B) = \frac{P(A_3|B)}{P(B)} = \frac{P(A_3) \cdot C_4^1 \left(\frac{2}{3}\right)^1 \left(\frac{1}{3}\right)^3}{P(B)} = \frac{64}{209}$$

记 $C = \{\text{下次还可以获得回报}\}$

$$P(C|B) = P(A_1|B)P((C|B)|A_1|B) + P(A_2|B)P((C|B)|A_2|B) + P(A_3|B)P((C|B)|A_3|B)$$

显然 $P(C|B)|A_1|B) = P(C|A_1) = \frac{1}{3}$, $P(C|B)|A_2|B) = P(C|A_2) = \frac{1}{3}$, $P(C|B)|A_3|B) = P(C|A_3) = \frac{1}{3}$

$$\therefore P(C|B) = \frac{1}{3} \cdot \frac{64}{209} + \frac{1}{3} \cdot \frac{81}{209} + \frac{1}{3} \cdot \frac{64}{209} \\ = \frac{1}{2}$$

18. 记 $A_i = \{\text{他通过第}i\text{个考试的事件}\}$, $B = \{\text{获得最终资格}\}$

$$P(A_1) = p, \quad P(A_1^c) = 1-p$$

$$P(A_2|A_1) = p, \quad P(A_2|A_1^c) = \frac{p}{2}, \quad P(A_2^c|A_1) = 1-p, \quad P(A_2^c|A_1^c) = 1-\frac{p}{2}$$

$$P(A_2) = P(A_1)P(A_2|A_1) + P(A_1^c)P(A_2|A_1^c) = p \cdot p + \frac{1}{2}(1-p) \cdot \frac{p}{2} = \frac{p+p^2}{2}$$

$$P(A_2^c) = 1 - P(A_2) = \frac{2-p-p^2}{2}$$

$$P(A_3|A_2) = p, \quad P(A_3|A_2^c) = \frac{p}{2}$$

$$P(A_3) = P(A_2)P(A_3|A_2) + P(A_2^c)P(A_3|A_2^c) = \frac{p+p^2}{2} \cdot p + \frac{1}{2} \frac{2-p-p^2}{2} \cdot \frac{p}{2} = \frac{2p+p^2+p^3}{4}$$

$$P(A_3^c) = \frac{4-2p-p^2-p^3}{4}$$

$$P(A_4) = P(A_3)P(A_4|A_3) + P(A_3^c)P(A_4|A_3^c) = \frac{2p+p^2+p^3}{4} \cdot p + \frac{4-2p-p^2-p^3}{4} \cdot \frac{p}{2} = \frac{4p+2p^2+p^3+p^4}{8}$$

$$P(A_4^c) = \frac{8-4p-2p^2-p^3-p^4}{8}$$

但是上面这些并没有用

$$P(B) = P(A_1 A_2 A_3 A_4) + P(A_1 A_2 A_3 A_4^c) + P(A_1 A_2 A_3^c A_4) + P(A_1 A_3^c A_2 A_4) + P(A_1^c A_2 A_3 A_4)$$



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$$\begin{aligned}
 P(A_1 A_2 A_3 A_4) &= P(A_1) P(A_2 | A_1) P(A_3 | A_1 A_2) P(A_4 | A_1 A_2 A_3) \\
 &= P(A_1) P(A_2 | A_1) P(A_3 | A_1) P(A_4 | A_1) \\
 &= p \cdot p \cdot p \cdot p = p^4 \\
 P(A_1 A_2 A_3 A_4^c) &= P(A_1) P(A_2 | A_1) P(A_3 | A_1) P(A_4^c | A_1) = p \cdot p \cdot p \cdot \frac{1}{2} = \frac{1}{2} (1-p)p^3 \\
 P(A_1 A_2^c A_3 A_4) &= P(A_1) P(A_2^c | A_1) P(A_3 | A_1) P(A_4 | A_1) = p \cdot (1-p) \cdot \frac{1}{2} \cdot p = (1-p) \frac{p^2}{2} \\
 P(A_1^c A_2 A_3 A_4) &= P(A_1^c) P(A_2 | A_1^c) P(A_3 | A_1^c) P(A_4 | A_1^c) = (1-p) \cdot \frac{1}{2} \cdot p \cdot p = (1-p) \frac{p^2}{2} \\
 P(B) &= p^4 + p^3 - p^4 + \frac{3}{2}(p^2 p^4) = \frac{5p^3}{2} - \frac{3p^4}{2} = \frac{5p^3 - 3p^4}{2}
 \end{aligned}$$

23. (1) $P(Y_1 \geq 0) = P(Y_1 = 0) + \frac{1}{2}$ if $A = \{x \in \mathbb{R}^4 : 1 \leq n \leq 4, Y_n \geq 0\}$

$$P(Y_1 \geq 0) = P(Y_1 = 0) + P(Y_1 = 1) \quad P(A) = P(Y_1 \geq 0) P(Y_2 \geq 0 | Y_1 \geq 0) P(Y_3 \geq 0 | Y_1, Y_2 \geq 0) P(Y_4 \geq 0 | Y_1, Y_2, Y_3 \geq 0)$$

$$\underline{P(Y_n \geq 0)} =$$

$$P(Y_1 \geq 0) = \frac{1}{2}$$

$Y_1 \geq 0$ if $Y_1 = 1$ 的情况下, 无论如何走 $Y_2 = 0$ 或 2 , $Y_2 \geq 0$

$$P(Y_2 \geq 0 | Y_1 \geq 0) = 1$$

$Y_2 = 0$ 的情况下, $P(Y_2 = 0) = P(Y_2 = 2) = \frac{1}{2}$, 仅当 $Y_2 = 0$ 且向左走时, $Y_3 < 0$

$$\text{故 } P(Y_3 \geq 0 | Y_1, Y_2 \geq 0) = \frac{1}{2} \times 1 + \frac{1}{2} \times \frac{1}{2} = \frac{3}{4}$$

$Y_3 \geq 0$ 的情况下, $Y_3 = 1$ 或 3 , 无论如何走 $Y_4 \geq 0$ 恒成立

$$\therefore P(A) = \frac{1}{2} \cdot 1 \cdot \frac{3}{4} \cdot 1 = \frac{3}{8}$$

(2) If $B = \{x \in \mathbb{R}^4 : 1 \leq n \leq 4, |Y_n| \leq 2\}$

$$P(B) = P(|Y_1| \leq 2) P(|Y_2| \leq 2 | |Y_1| \leq 2) P(|Y_3| \leq 2 | |Y_1|, |Y_2| \leq 2) P(|Y_4| \leq 2 | |Y_1|, |Y_2|, |Y_3| \leq 2)$$

$$P(|Y_1| \leq 2) \text{ 与 } P(|Y_2| \leq 2 | |Y_1| \leq 2) \text{ 显然为 } 1$$

$$P(|Y_3| \leq 2 | |Y_1| \leq 2, |Y_2| \leq 2) = P(Y_3 = 0) + \frac{1}{2} P(Y_3 = 2) + \frac{1}{2} P(Y_3 = -2) = \frac{3}{4}$$

在 $|Y_1| \leq 2, |Y_2| \leq 2$ 的前提下, $|Y_3| \leq 2$ 当且仅当 $Y_3 = -1$ 或 0 或 1

故 $|Y_4| \leq 2$ 恒成立, $P(|Y_4| \leq 2 | |Y_1|, |Y_2|, |Y_3| \leq 2) = 1$

$$P(B) = \frac{3}{4}$$



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(3) 记 $C = \{x \in \mathbb{R}^n : 1 \leq n \leq 4, Y_n \geq 0\}$

$$P(A | Y_4=0) = \frac{P(A \cap Y_4=0)}{P(Y_4=0)} = \frac{P(A) P(Y_4=0 | A)}{P(Y_4=0)}$$

$$P(Y_4=0) = \binom{4}{4} \left(\frac{1}{2}\right)^4 = \frac{1}{16}$$

$$\text{故 } P(Y_3=1 | Y_1, Y_2 \geq 0) = \frac{2}{3}, \quad P(Y_3=3 | Y_1, Y_2 \geq 0) = \frac{1}{3}$$

$$\text{故 } P(Y_4=0 | A) = \frac{1}{2} P(Y_3=1 | Y_1, Y_2 \geq 0) = \frac{1}{3}$$

$$P(A | Y_4=0) = \frac{\frac{1}{3} \cdot \frac{1}{16}}{\frac{1}{16}} = \frac{1}{3}$$

25. (1) 记得到正面的次数为 X , $X = 0, 1, 2, 3, 4, 5, 6$

$$\begin{aligned} P(X=k) &= \frac{1}{6} \binom{k}{k} \frac{1}{2^k} + \frac{1}{6} \binom{k+1}{k} \frac{1}{2^{k+1}} + \dots + \frac{1}{6} \binom{6}{k} \frac{1}{2^6} \\ &= \frac{1}{6} \sum_{n=k}^6 \binom{n}{k} \frac{1}{2^n} \end{aligned}$$

(2) 显然 $n=3, 4, 5, 6$, 记 Y 为得到的点数

$$P(Y=n | X=3) = \frac{P(Y=n, X=3)}{P(X=3)} = \frac{P(X=3 | Y=n) P(Y=n)}{P(X=3)}$$

$$P(X=3) = \frac{1}{6} \sum_{n=3}^6 \binom{n}{3} \frac{1}{2^n} = \frac{1}{6} \left(\frac{1}{2^3} + 3 \cdot \frac{1}{2^4} + 10 \cdot \frac{1}{2^5} + 20 \cdot \frac{1}{2^6} \right) = \frac{5}{32}$$

$$P(Y=n) = \frac{1}{6} \quad P(X=3 | Y=n) = \binom{n}{3} \frac{1}{2^3} \cdot \frac{1}{2^{n-3}} = \frac{\binom{n}{3}}{2^n}$$

$$\text{故 } P(Y=n | X=3) = \frac{\frac{\binom{n}{3}}{2^n}}{\frac{5}{32}} = \frac{16}{15} \cdot \frac{\binom{n}{3}}{2^n}$$

28. ~~$P(X \leq Y) = P(Y=1) P(X \leq Y | Y=1) + P(Y=2) P(X \leq Y | Y=2)$~~



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28. 将 Ω 分为不可数个集合 $\{Y=1\}, \{Y=2\}, \dots, \{Y=n\}, \dots$

$$\begin{aligned} P(X \leq Y) &= P(Y=1)P(X \leq Y | Y=1) + P(Y=2)P(X \leq Y | Y=2) + \dots + P(Y=n)P(X \leq Y | Y=n) + \dots \\ &= \sum_{k=1}^{\infty} P(Y=k)P(X \leq k) \\ &= \sum_{k=1}^{\infty} \left[P(Y=k) \sum_{i=1}^k P(X=i) \right] \\ &= \sum_{k=1}^{\infty} \left(p_k \sum_{i=1}^k p_i \right) \end{aligned}$$

$$\begin{aligned} P(X=Y) &= P(Y=1)P(X=Y | Y=1) + P(Y=2)P(X=Y | Y=2) + \dots + P(Y=n)P(X=Y | Y=n) + \dots \\ &= \sum_{k=1}^{\infty} P(Y=k)P(X=k) \\ &= \sum_{k=1}^{\infty} p_k^2 \end{aligned}$$

41. 先证 $\varphi(s+t) = \varphi(s)\varphi(t) \Leftrightarrow \varphi(t) = e^{-\lambda t} \quad (s, t \geq 0)$
 $\Rightarrow s=t=0, \varphi(0)=\varphi(t), \varphi(0)=1 \quad (\varphi(0)=0 \text{ 不成立})$
 故 $\varphi(\frac{1}{n}) = \alpha^{\frac{1}{n}}$; $\varphi(\frac{m}{n}) = \varphi(\frac{m-1}{n})\varphi(\frac{1}{n}) = \varphi(\frac{m-2}{n})\varphi(\frac{1}{n}) = \dots = \varphi^{m-1}(\frac{1}{n}), n \in \mathbb{N}^+$
 故 $\varphi(t) = \alpha^t$ 在有理数上成立
 对 $\forall r \in \mathbb{R}^+ \setminus \mathbb{Q}^+$, 取 $\{a_n\}, \{b_n\}$ 使得 $a_n \leq b_n$ 且 $\lim_{n \rightarrow \infty} a_n = \lim_{n \rightarrow \infty} b_n = r, a_n, b_n \in \mathbb{Q}^+$
 因 φ 单调不增则 $\varphi(b_n) \leq \varphi(a_n) \leq \varphi(r)$, 令 $n \rightarrow \infty$ 并由夹逼原理

$$\begin{aligned} \varphi(r) &= \lim_{n \rightarrow \infty} \varphi(a_n) = \lim_{n \rightarrow \infty} \alpha^{a_n} = \alpha^r, \quad r \text{ 是正无理数} \\ \text{故 } \varphi(t) &= \alpha^t = \varphi^{(1)} = \alpha^{rt} e^{-\lambda t}, \quad \text{其中 } \lambda = \ln \frac{1}{\varphi(1)} \geq 0 \end{aligned}$$

下证满足 $P(T > s+t | T > s) = P(T > t), (s, t \geq 0) \Leftrightarrow T$ 满足指数分布

$$\Rightarrow P(T > s+t | T > s) = \frac{P(T > s+t)}{P(T > s)} = P(T > t), \quad P(T > t) \text{ 关于 } t \text{ 显然单减非增}$$

$$\text{则 } P(T > t) = e^{-\lambda t} \quad t > 0$$

$$F_T(t) = 1 - P(T > t) = 1 - e^{-\lambda t} \quad (t > 0), \quad T \text{ 满足指数分布}$$

$$\Leftarrow T \text{ 满足 } F_T(t) = 1 - e^{-\lambda t}, \text{ 故 } P(T > t) = 1 - F_T(t) = e^{-\lambda t}$$

$$P(T > s+t) = e^{-\lambda(s+t)} = e^{-\lambda s} \cdot e^{-\lambda t} = P(T > s)P(T > t)$$

$$\therefore P(T > s+t | T > s) = P(T > s+t) / P(T > s) = P(T > t)$$

证毕



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